

Understanding the Internet

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An engineering approach

The Internet is an engineered artifact.

- Designed under constraints for a purpose.

Therefore to understand the Internet:

- Understand **constraints**.
- Understand **purposes**.
- Understand **solutions** (and create new ones)!

A scientific approach

The Internet is a **complex system**.

- Understanding components is NBNS.
- Understanding interactions among components is NBNS.

Understanding requires:

- Appropriate abstraction.
- Observation of **emergent properties**.

Traffic jams

Example: automobile traffic is a complex system.

What is a traffic jam?

- Made of cars, but not always the same cars.
- It's an abstract entity, like a storm front.
- It has properties that its components don't have.

Traffic jams move **backwards**.

Outline

Exercise these viewpoints:

- The bottom-up Internet.
- The emergent Internet.

Physical layer

Starting with a medium that transmits an **analog signal**...



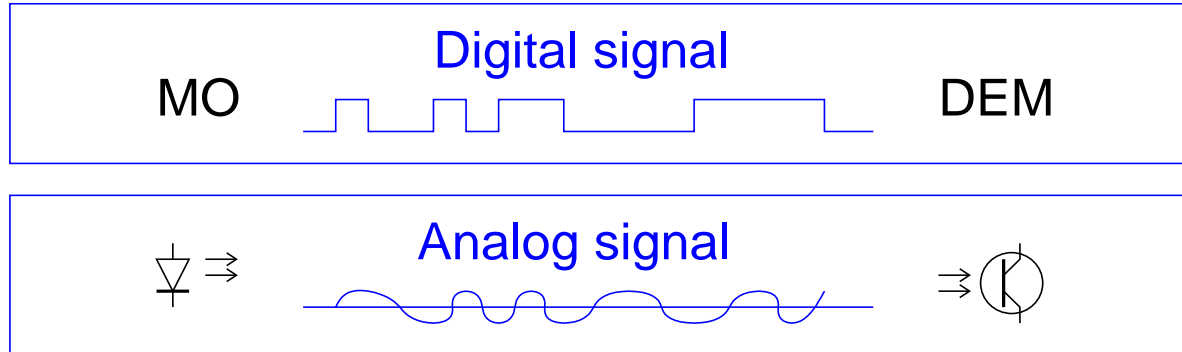
...how can we transmit a **digital signal**?

Solution: modulation.

- Vary amplitude, frequency, etc.
- Distinguish (at least) two levels.

Signal layer

Physical layer + mo/dem = Signal layer



Abstractly, a modem transmits a **digital signal**.

Physically, there's no such thing!

HI and LO are **abstract symbols** with many possible physical representations.

Transmitting bits

- A bit can have two states, 0 or 1.
- A digital signal has two states, LO and HI.
- Naive encoding: LO represents 0, HI represents 1.

Problem?

Transmitting bits

- A bit can have two states, 0 or 1.
- A digital signal has two states, LO and HI.
- Naive encoding: LO represents 0, HI represents 1.

Problem?

1. Clock recovery.
2. Threshold discovery.

Manchester encoding

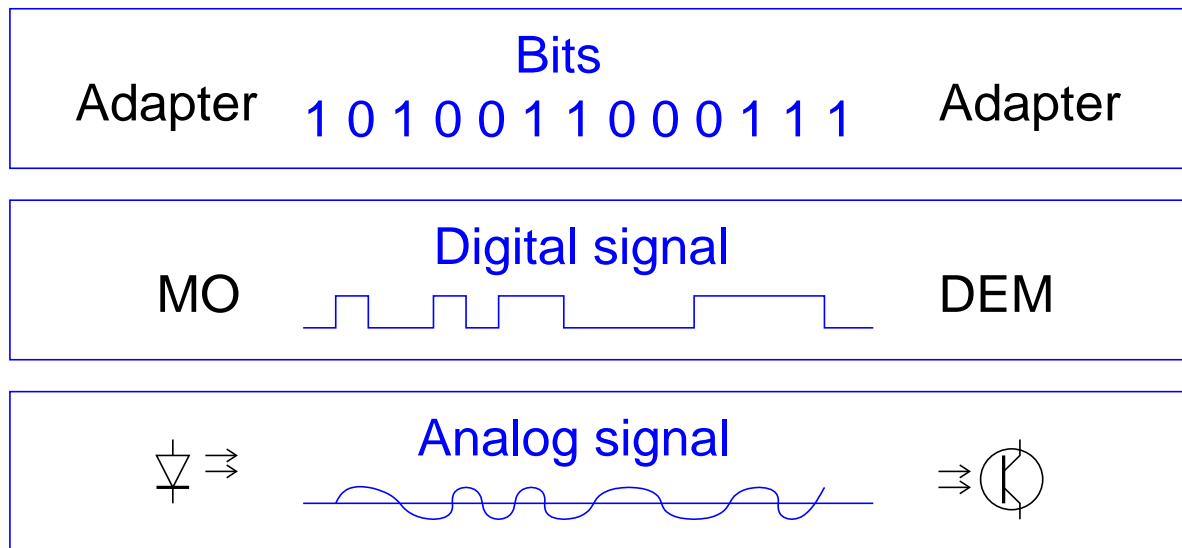
LO-HI represents 0.

HI-LO represents 1.

- Lots of transitions; good for clock recovery.
- Equal number of LO and HI; good for threshold discovery.

Bit layer

Signal layer + encoding = **Bit layer**.



Abstractly, an **adapter transmits bits**.

Physically, there's no such thing!

0 and 1 are abstract symbols with many possible physical representations.

Error detection

How do you know the bits you got are right?

- Group bits into **frames**.
- Attach **redundant information** to each frame.

How do we identify the beginning and end of a frame?

Framing

A frame consists of:



- Special sequence marks **beginning** and **end**.
- After the data, send a checksum.
- Checksum is a function of the data bits.

Problem?

Bit stuffing

What if **0111110** appears in the data?

Sender:

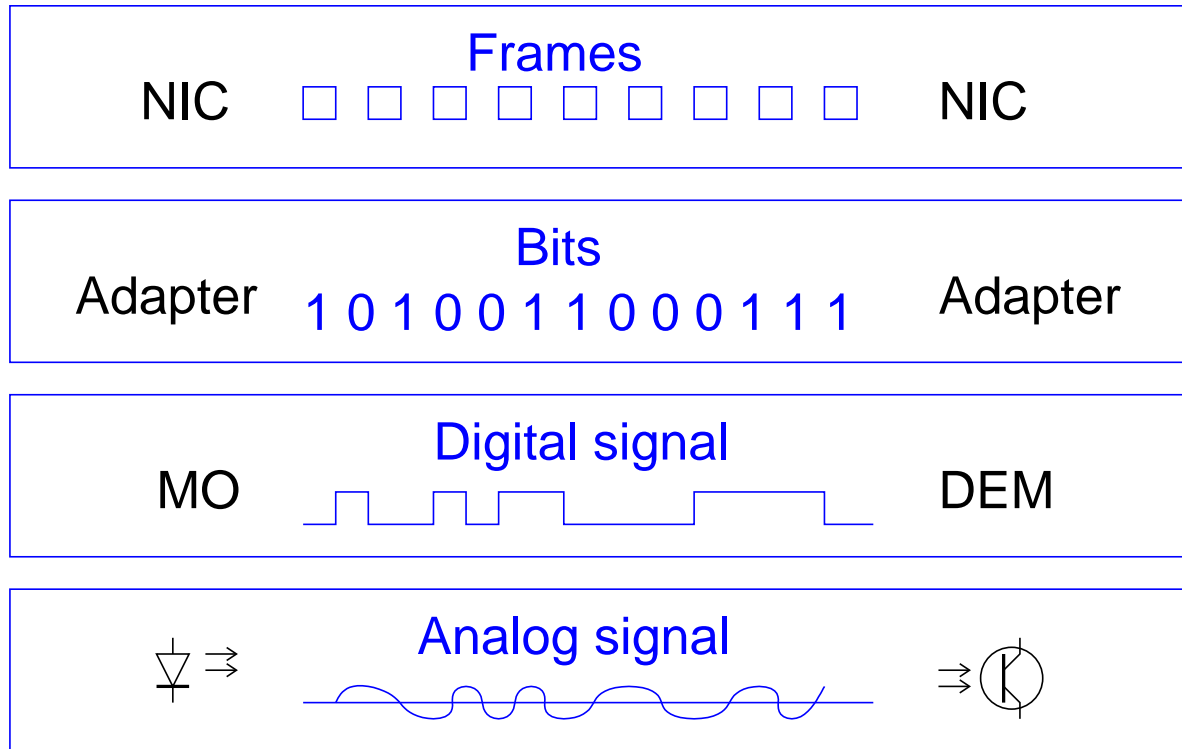
- Whenever you send **011111**, stuff in a **0**.

Receiver:

- Whenever you see **011111**, check the next bit:
 - If **0**, remove it.
 - If **1**, check the next bit:
 - If **0**, that's the end of frame.
 - If **1**, there must have been an error!

Link layer

Bit layer + framing + error detection = **Link layer**.



Abstractly, a **NIC transmits frames**.

Physically, there's no such thing!

The Internet Protocol

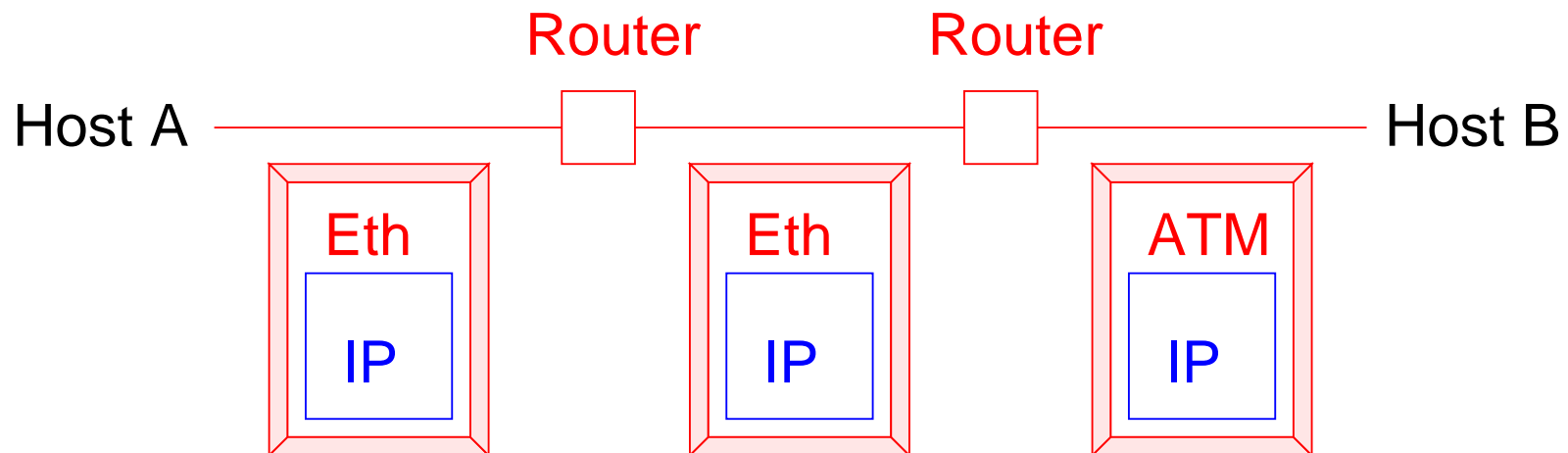
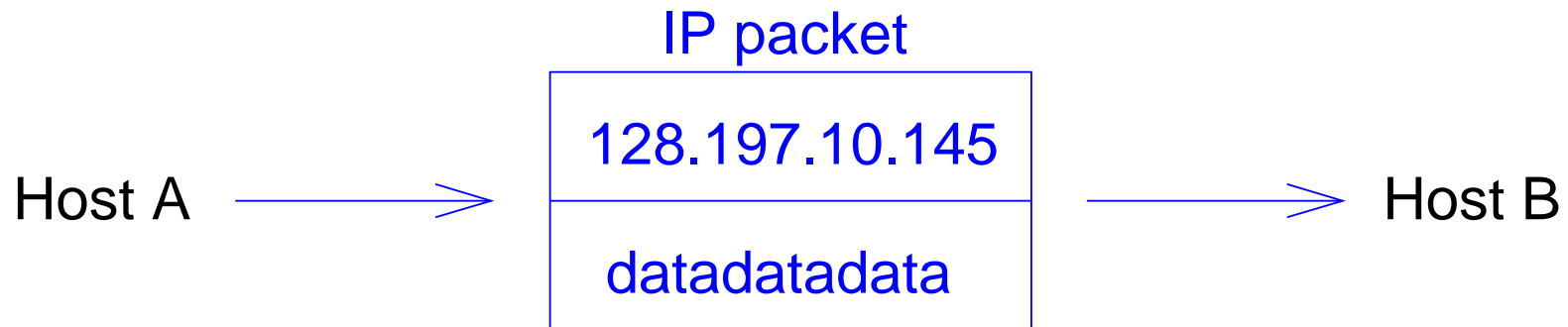
Links are **point-to-point**.

How do you talk to someone you're not connected to?

- Assign IP addresses (128.197.10.145).
- Relay **packets** through routers.

What's a packet?

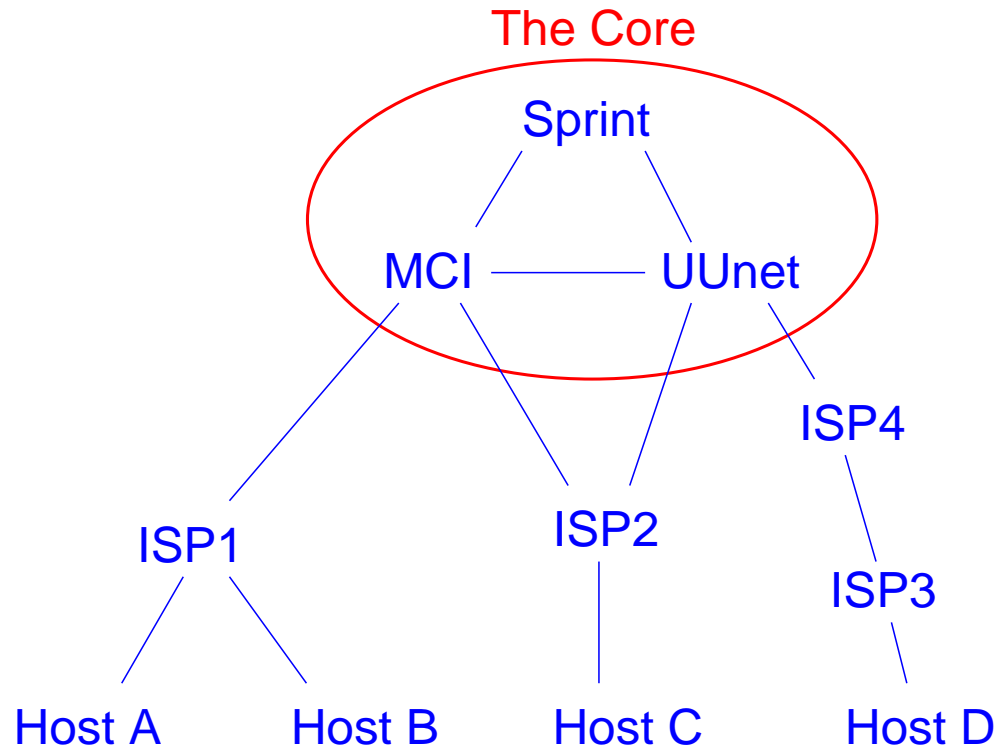
Abstractly, Host A transmits a packet to Host B.



Physically, the packet's data are copied from link to link.

Internet structure

Physical structure is mostly hierarchical.



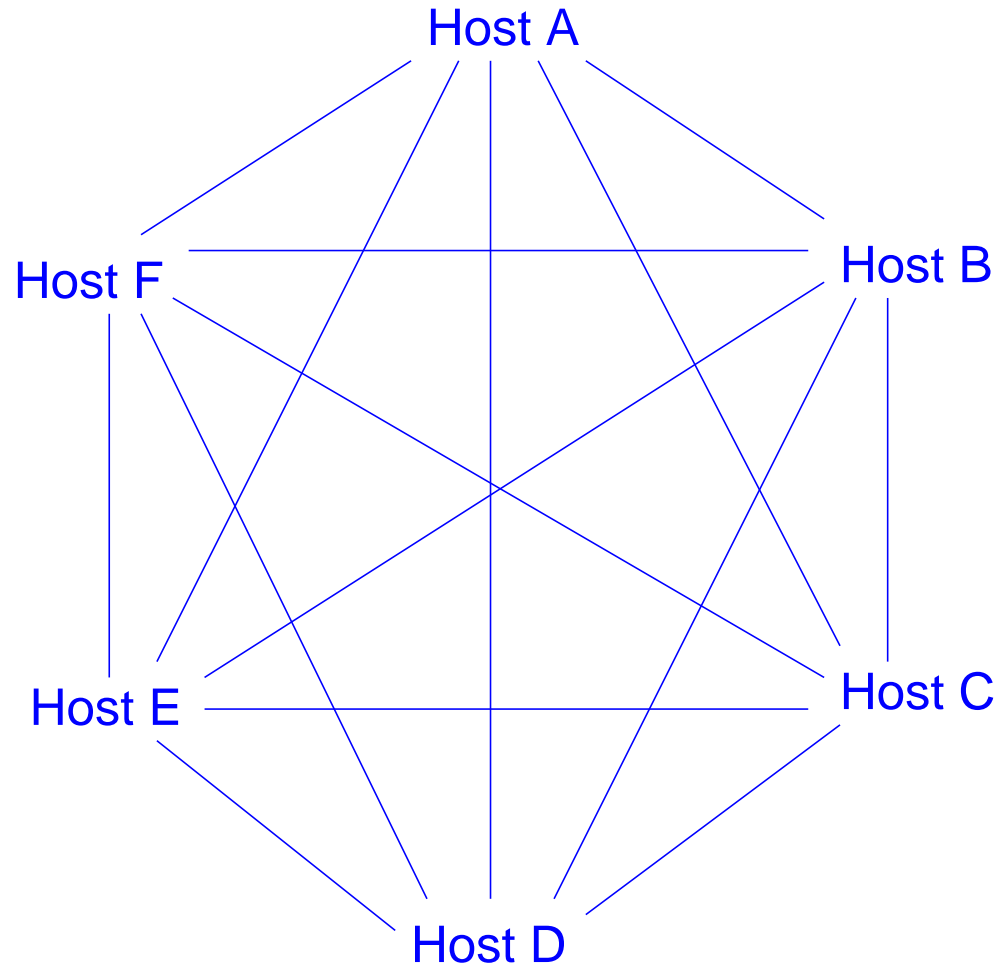
- Core providers completely connected.

A typical path

- 0 rosco.bu.edu (128.197.10.145)
- 1 **csnet**-gw (128.197.10.1)
- 2 extgw-dgw-**cumm111** (128.197.254.122)
- 3 **bos**-edge-02.inet.**qwest**.net (63.145.1.77)
- 4 **bos**-core-02.inet.**qwest**.net (205.171.28.29)
- 5 **jfk**-core-01.inet.**qwest**.net (205.171.8.19)
- 6 **jfk**-core-03.inet.**qwest**.net (205.171.230.6)
- 7 **jfk**-brdr-02.inet.**qwest**.net (205.171.230.25)
- 8 so-4-0-3.**nycmny1**-hcr3.**bbnplanet**.net (4.25.133.29)
- 9 p10-0.**nycmny1**-nbr2.**bbnplanet**.net (4.0.7.14)
- 10 so-4-0-0.**bstnma1**-nbr2.**bbnplanet**.net (4.24.6.49)
- 11 p2-0.**bstnma1**-cr10.**bbnplanet**.net (4.24.10.62)
- 12 multilink1.**folincollege2**.**bbnplanet**.net (4.24.90.218)
- 13 4.21.173.194 (4.21.173.194)
- 14 * * *
- 15 www.olin.edu (4.21.175.12)

Abstract structure

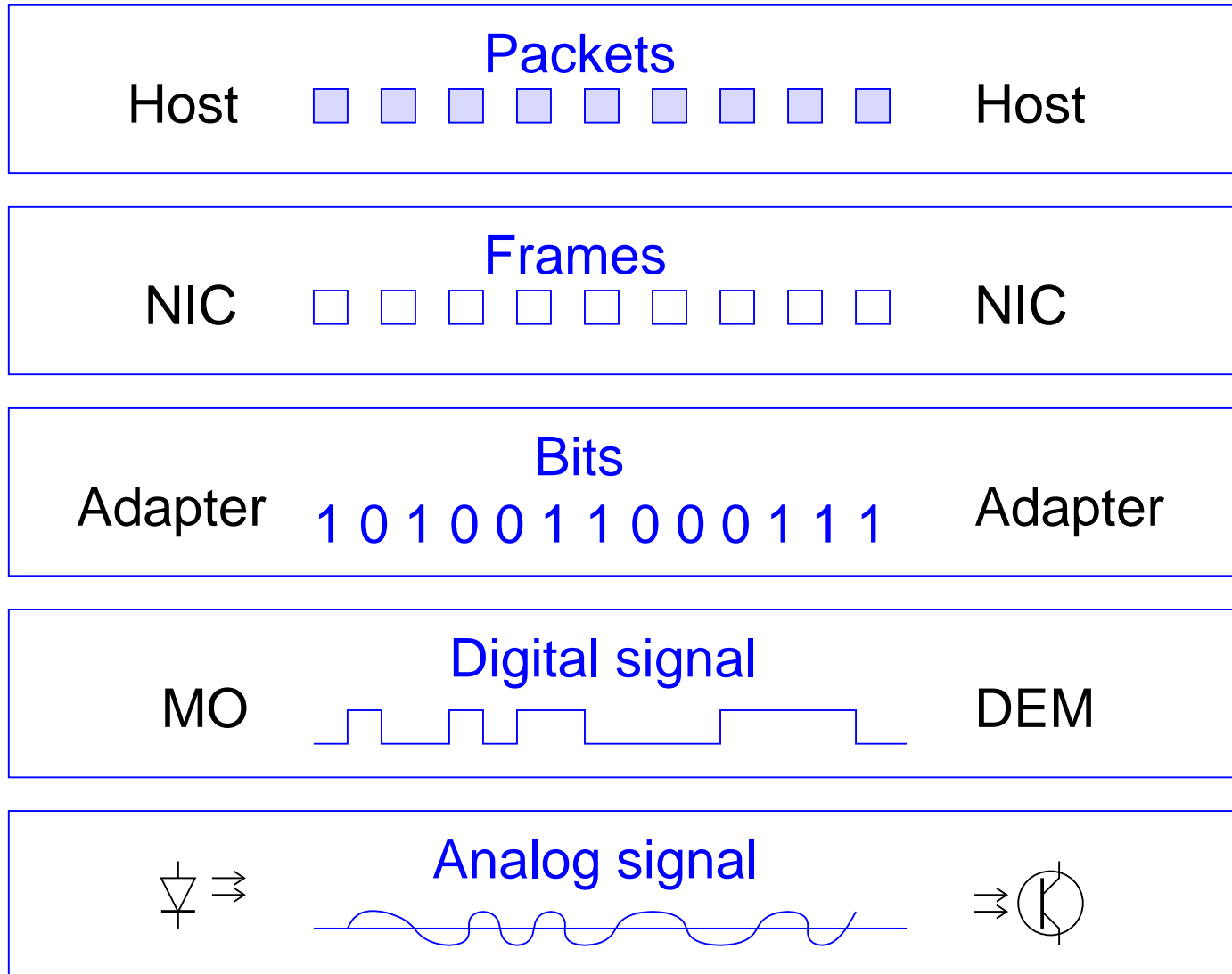
Fortunately, IP hides this complexity!



IP creates an abstract, **completely connected** network.

The Network Layer

Link layer + addressing + routing = Network Layer



Best effort delivery

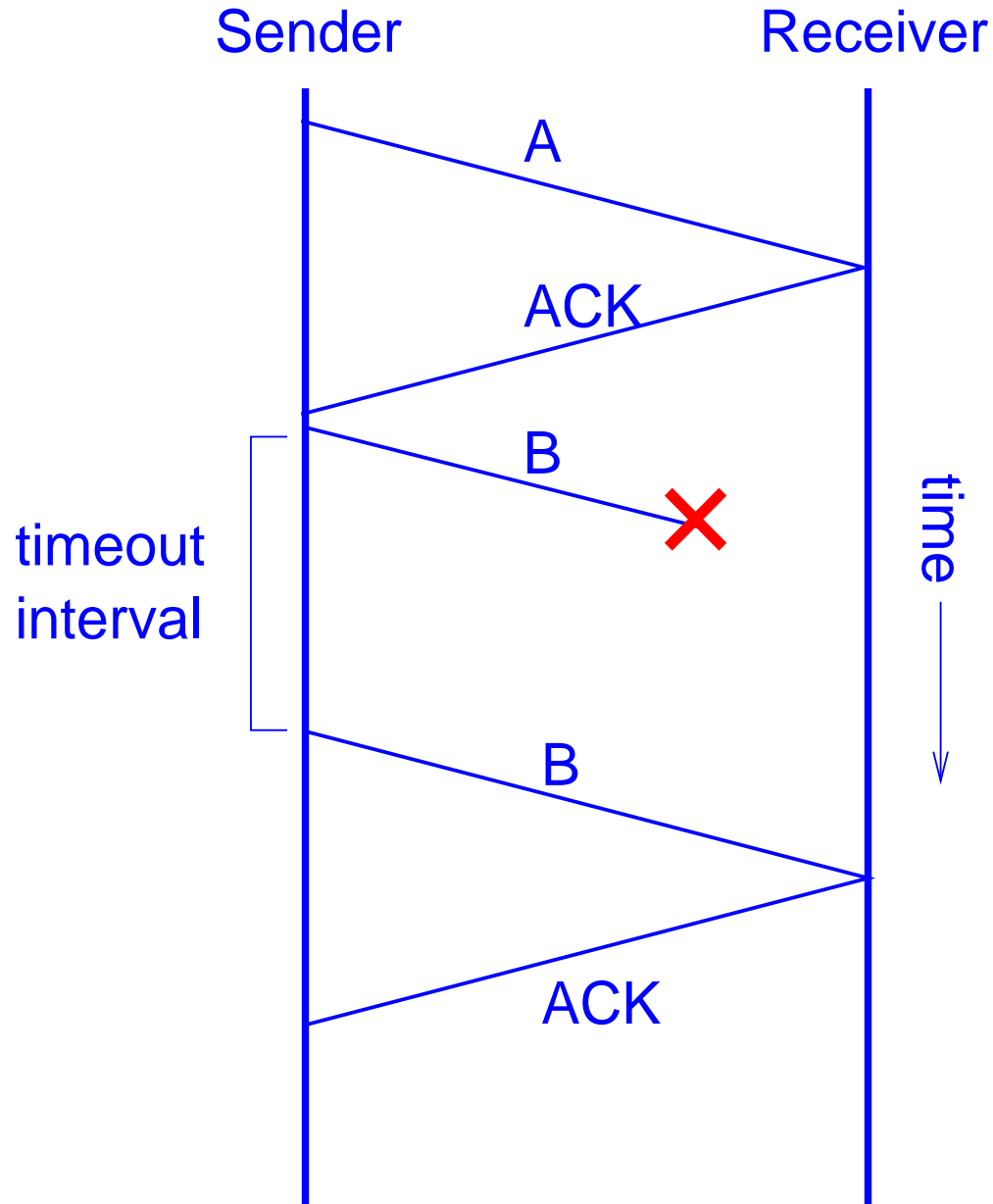
Absolutely, positively might get there eventually.

Who could ask for anything more?

- Exactly-once delivery.
- In-order delivery.

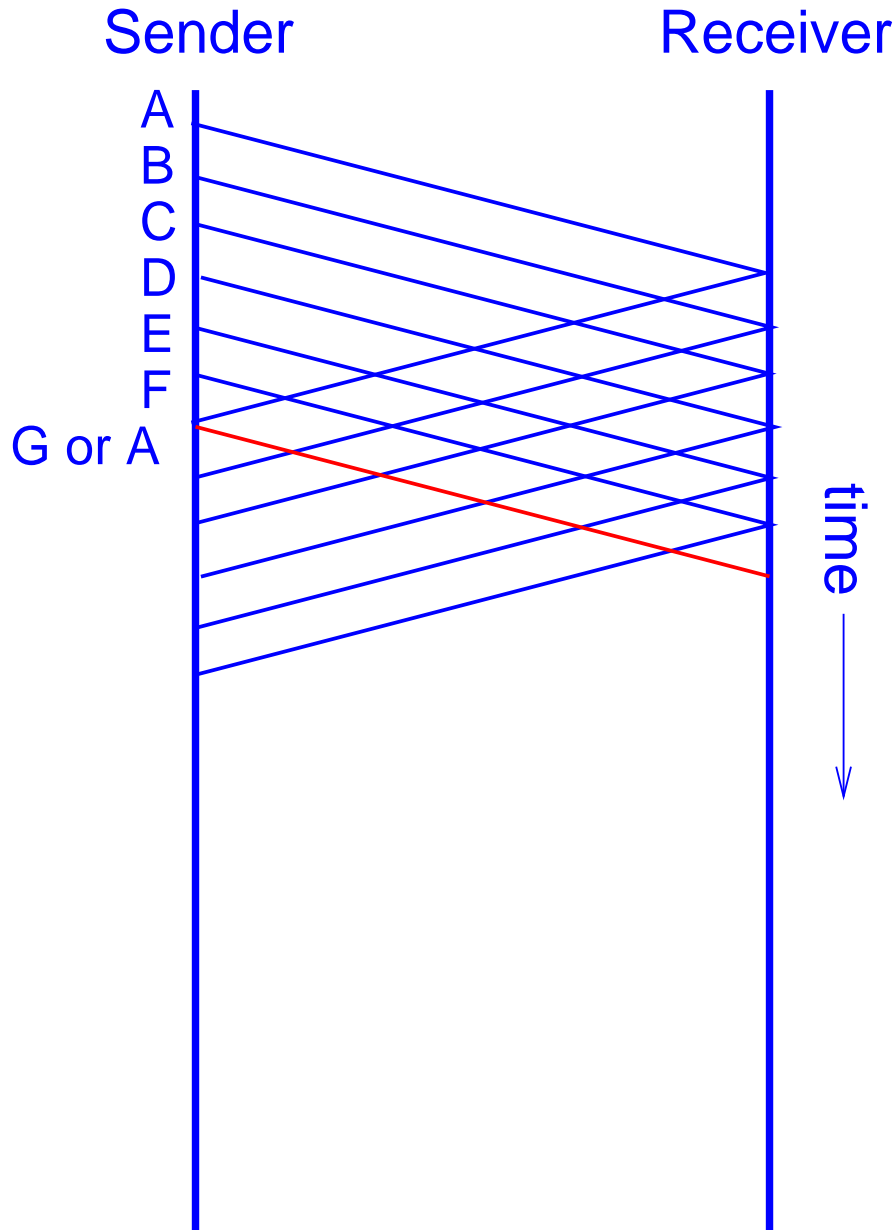
Can you build **reliable communication** on an **unreliable medium**?

Stop and wait



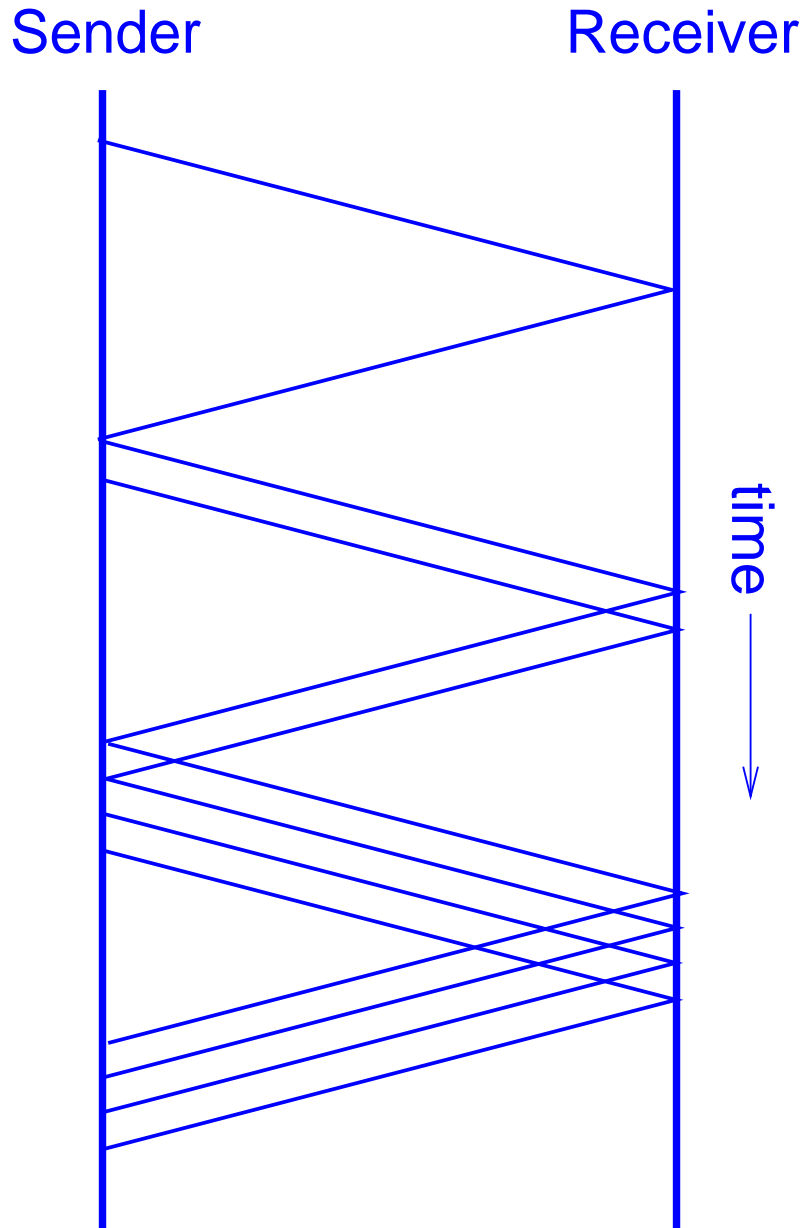
- Transmit.
- Acknowledge.
- Timeout... retransmit.
- Problem?

Sliding window



- Keep a copy of unACKed packets.
- Get an ACK, send a new packet.
- Miss an ACK, retransmit.
- Problem?

Slow start



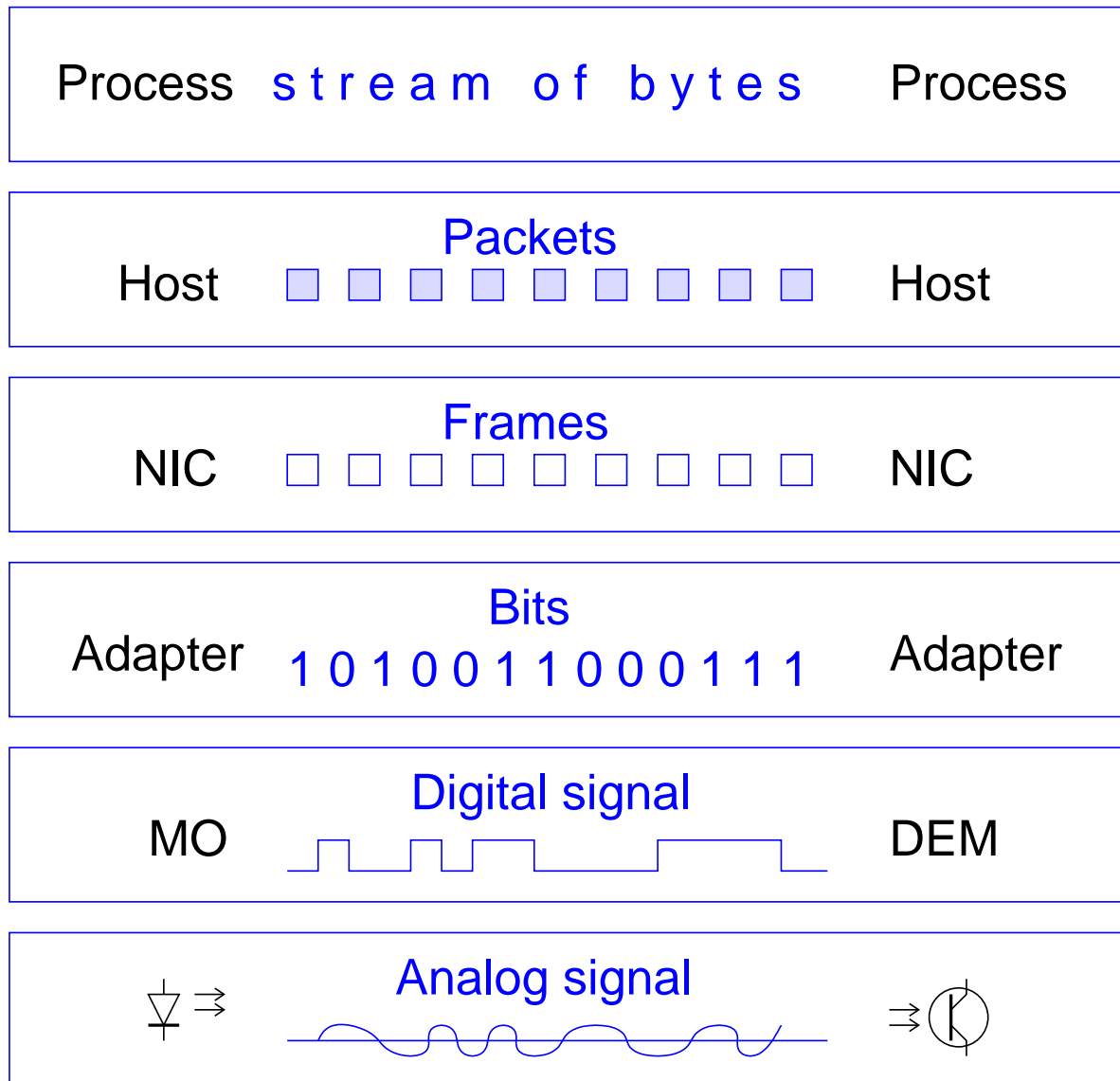
- Start with one packet.
- Get an ACK, send two packets.
- Watch for a **congestion signal**.

Transmission Control Protocol

- Reliable delivery (ACK-retransmit).
- Flow control (sliding window).
- Congestion control (slow start).
- Byte stream abstraction.
- Process-to-process communication.

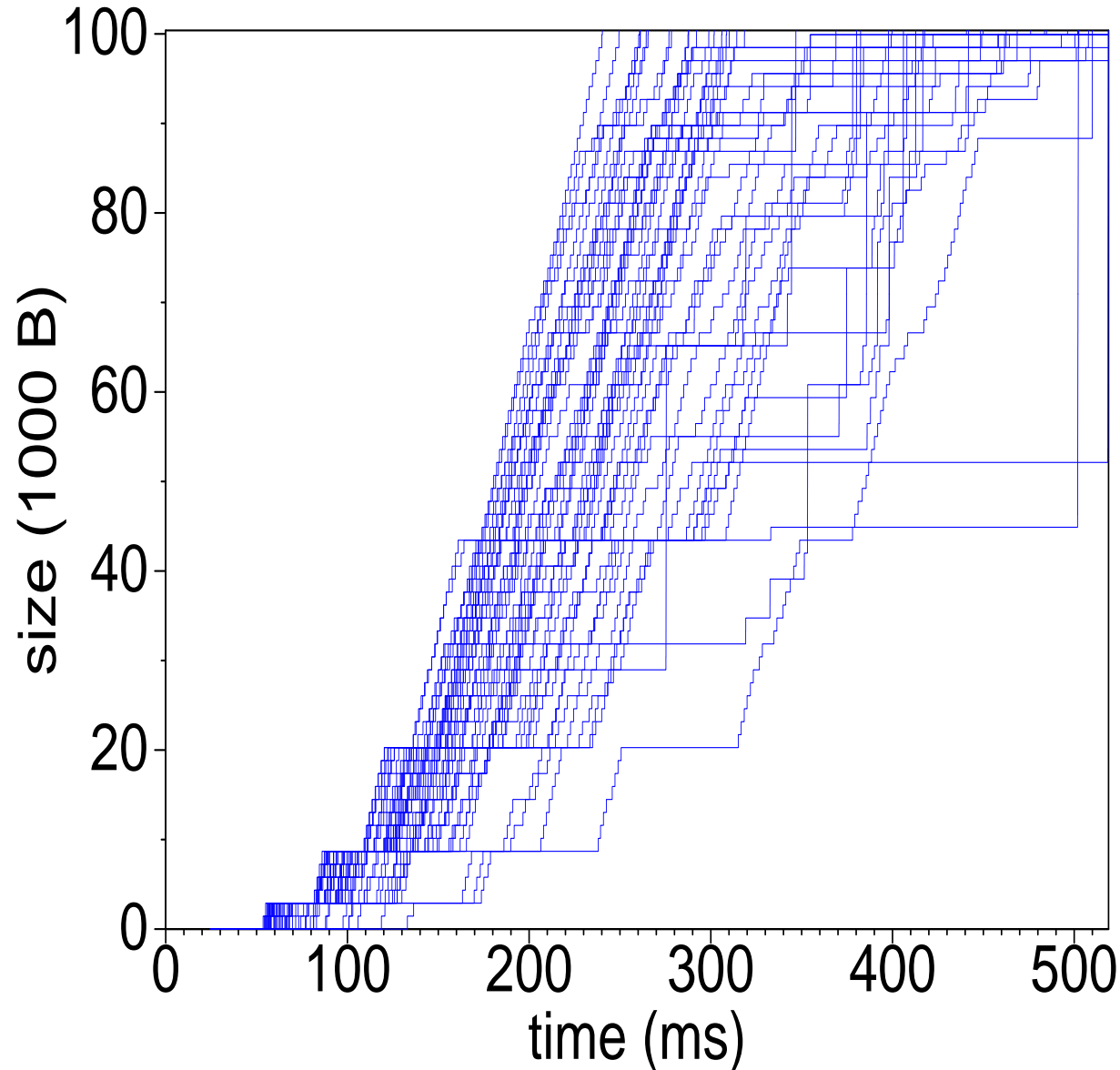
The Transport Layer

Network layer + sliding window + slow start = Transport layer



Modeling TCP Performance

server 7 timing chart

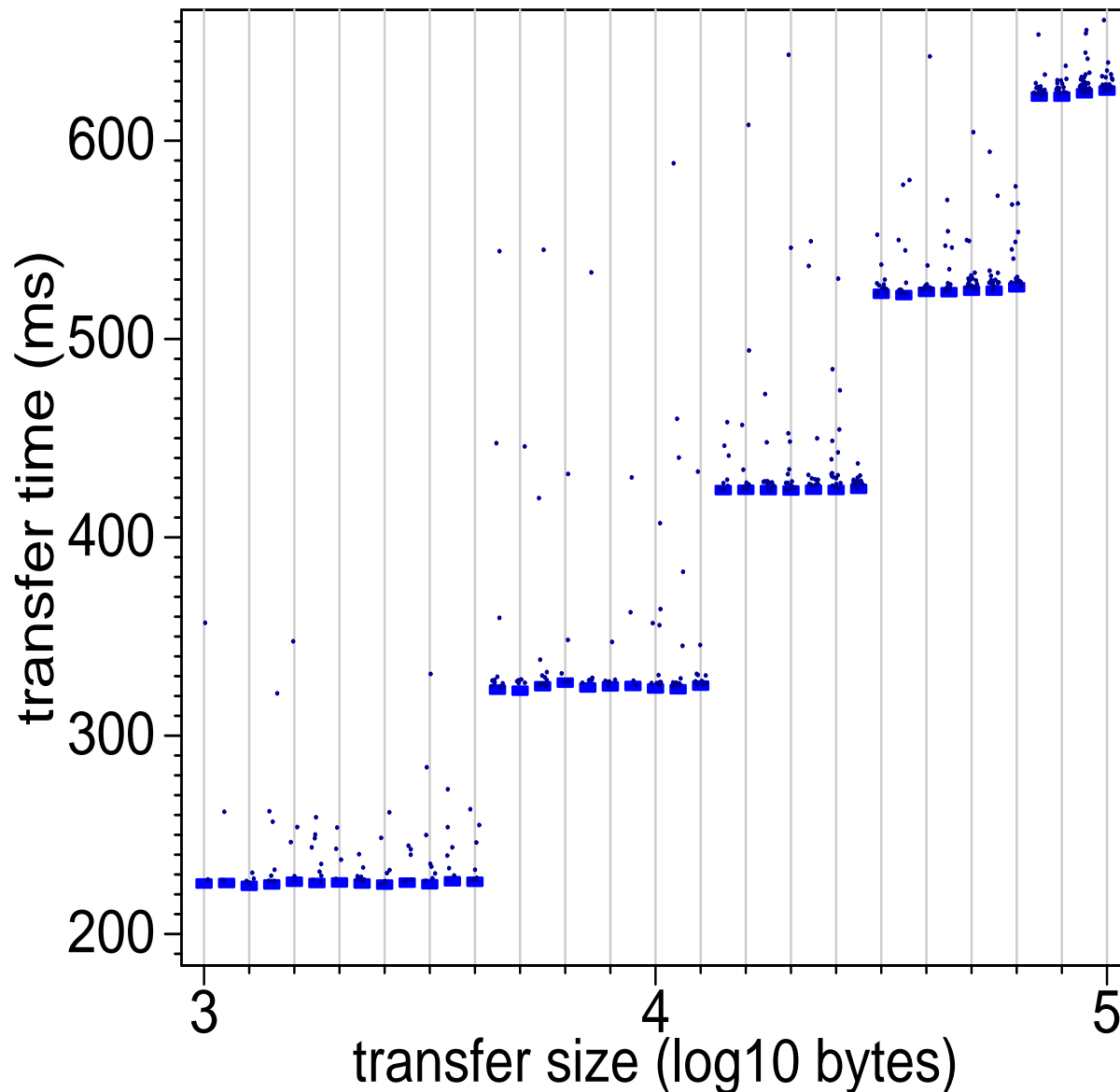


- Observe network transfers.
- Estimate path parameters.
- Predict transfer times.

“Modeling TCP Performance,”
Downey, 2003.

Modeling TCP Performance

HTTP transfer times



- Transfer times for short files are multiples of the **round trip time**.

Who uses TCP?

90% of Internet traffic is based on TCP:

- TCP + **HTTP** = The Web
- TCP + **SMTP** = email
- TCP + **gnutella protocol** = peer-to-peer network

Application protocols

A little HTTP:

```
GET index.html HTTP/1.1
```

```
HTTP/1.1 200 OK
```

```
Server: Apache-2.0.44
```

```
Content-Length: 12481
```

```
Content-Type: text/html
```

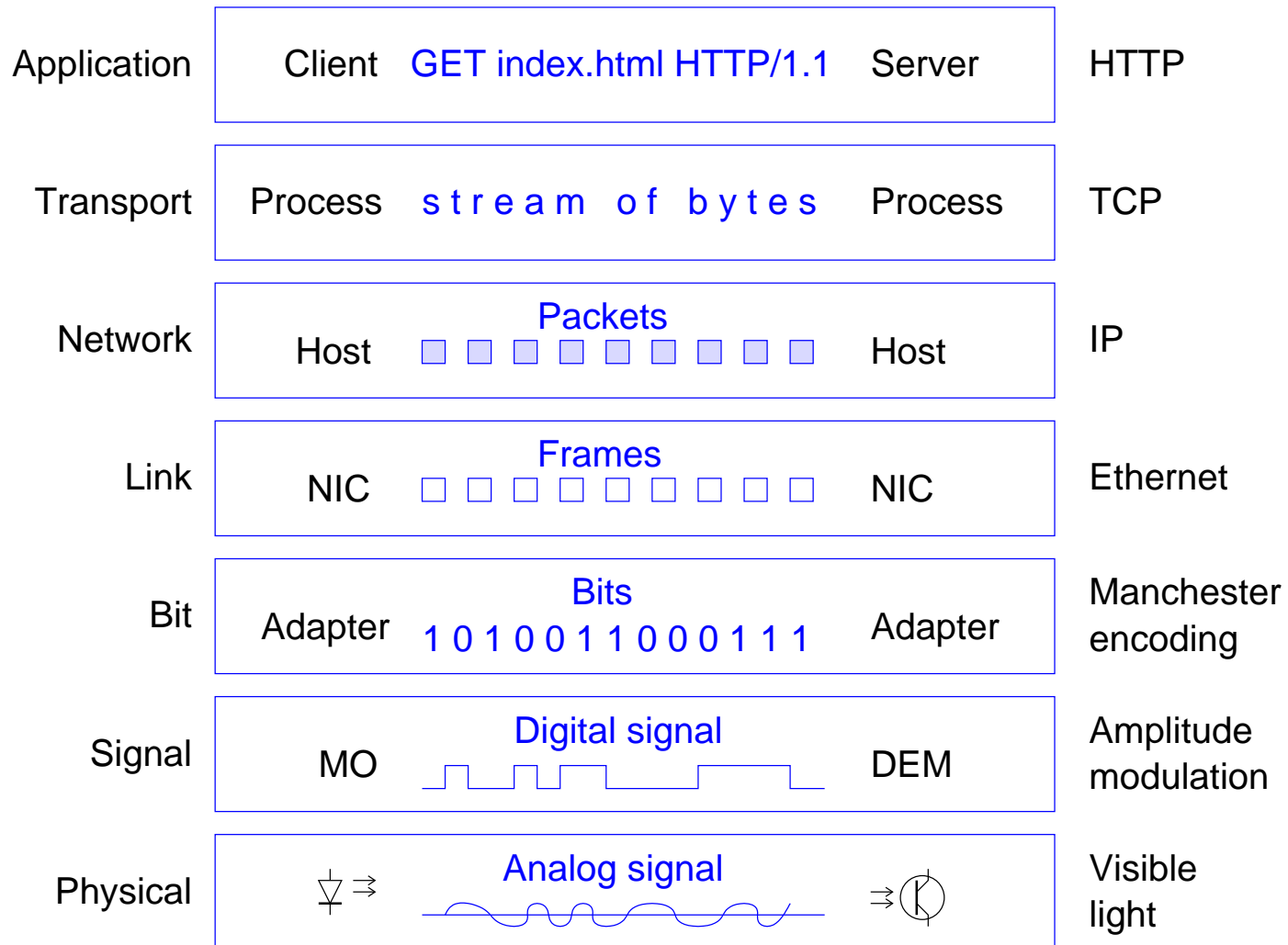
```
<html>
```

```
...
```

```
</html>
```

The Application Layer

Transport layer + application protocol = **Application layer**



So now you know...

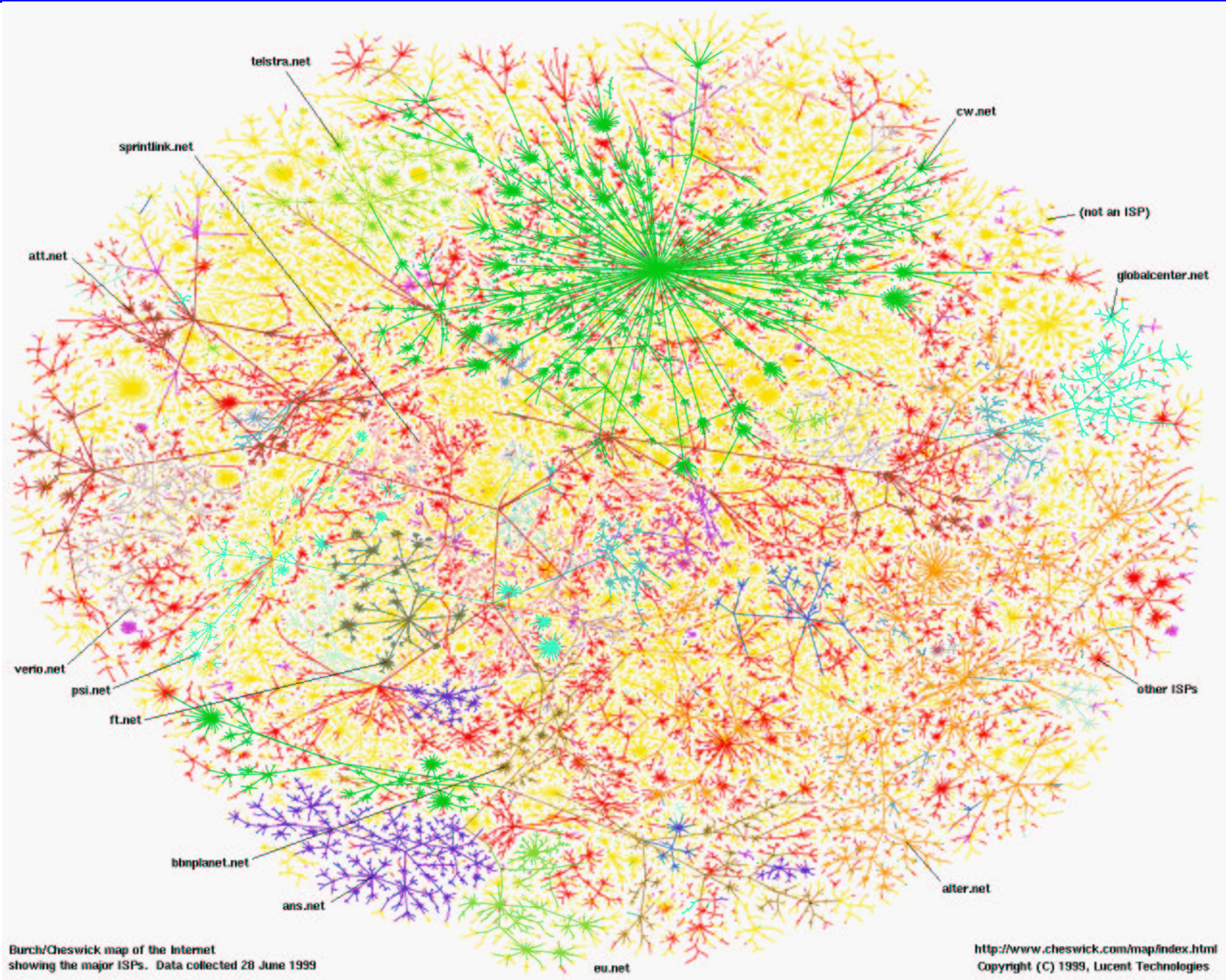
...given a medium that transmits an **analog signal**...



...how to build an **Internet-compatible host**.

But how do you build the Internet?

You don't. It grows...

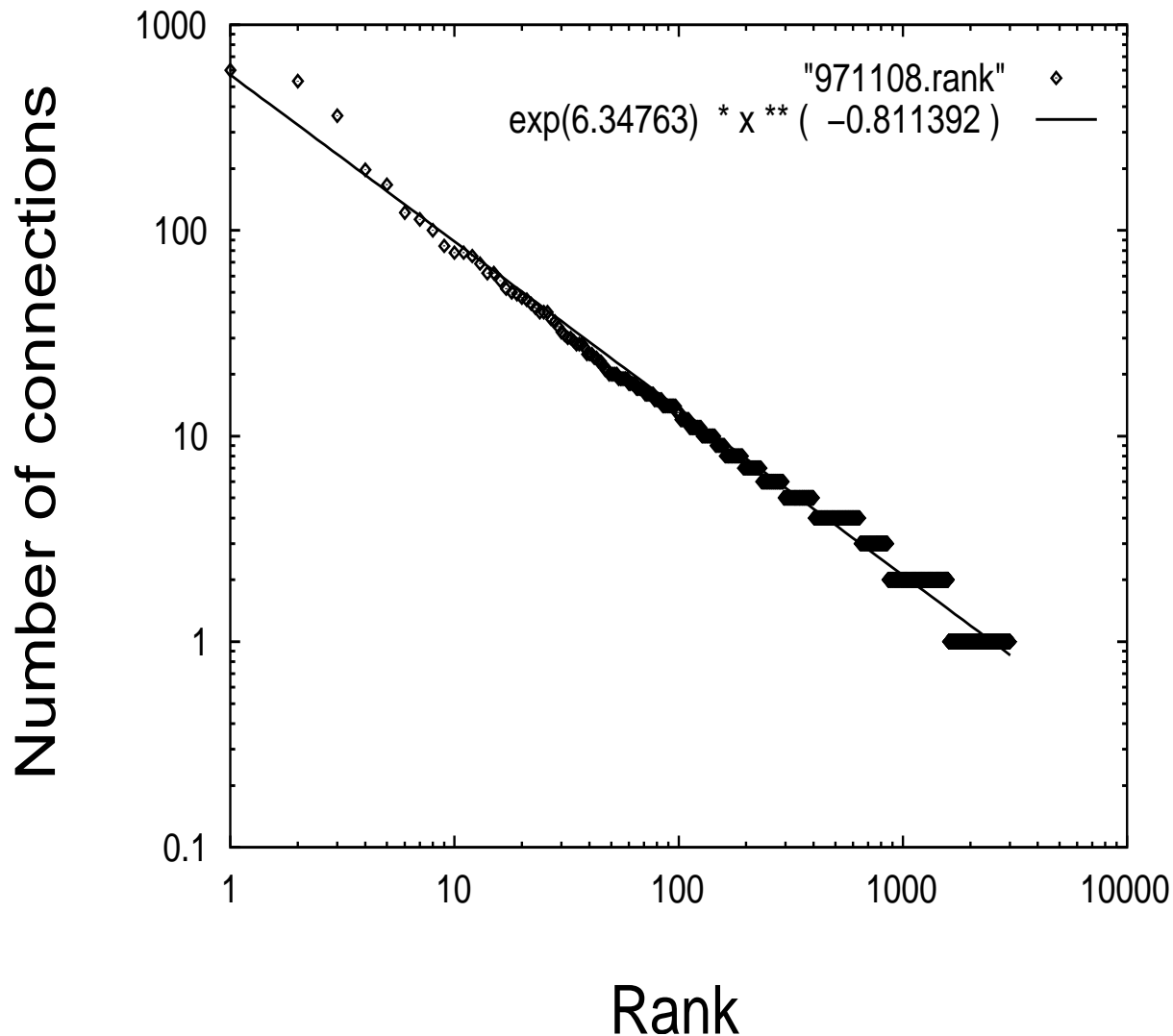


It grows, and grows...

Internet structure depends on:

- Technological details.
- Geographic details.
- Demographic details.
- Economic factors.
- Human factors.

It grows in mysterious ways



- Long-tailed distribution of connections (many small, some very large).
- Characteristic of **rich-get-richer** growth.

“On Power-Law Relationships of the Internet Topology” Faloutsos³, 1999.

Zipf's law

Zipf rank relationships in:

- Popularity of words (Zipf 1932).
- Sizes of cities (Zipf 1939).
- Individual income (Pareto 1906).
- Magnitude of earthquakes.
- Debt of bankrupt companies.

...and on and on.

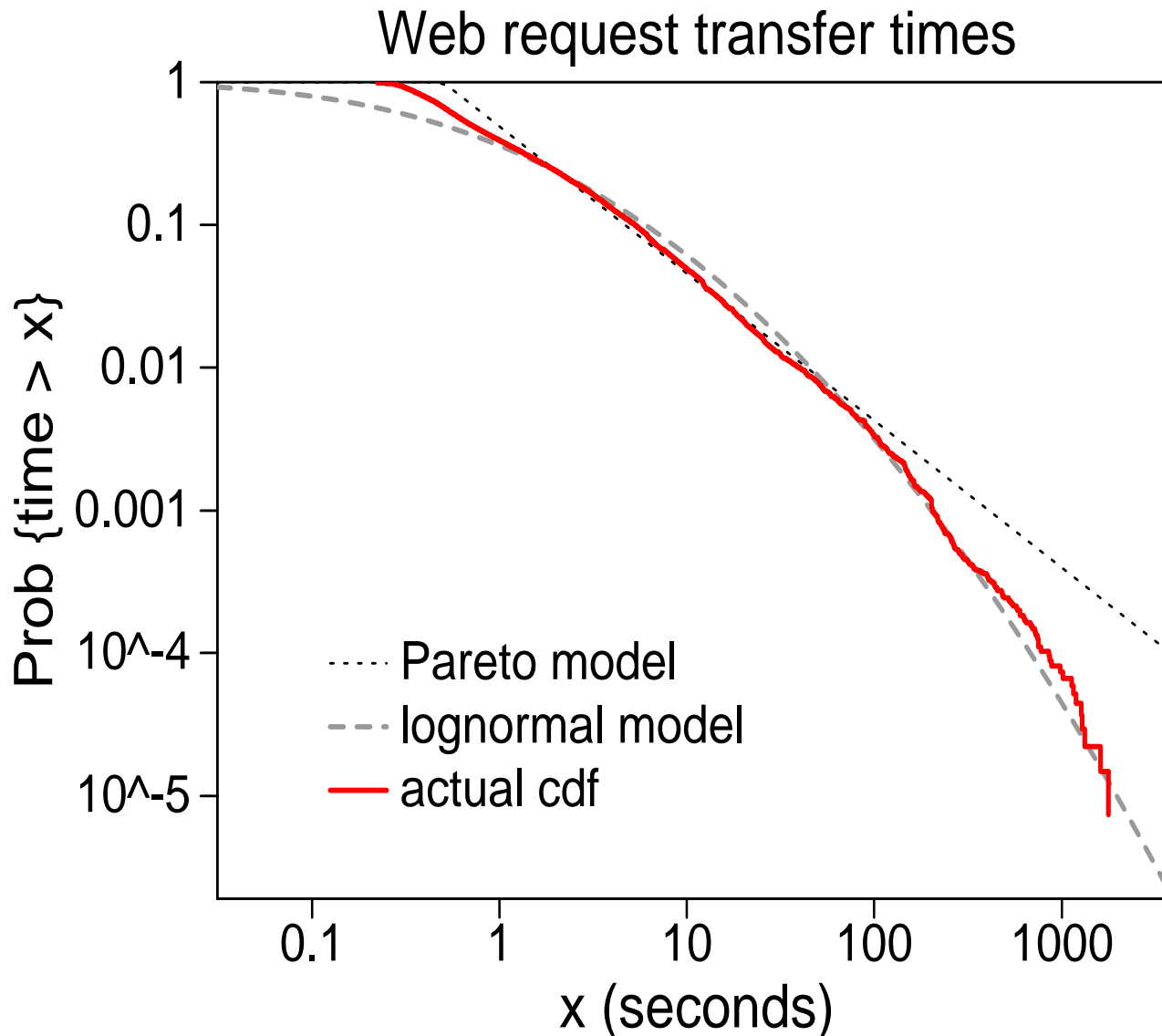
Zipf's law

And in the Internet...

- Degree of connectivity.
- Popularity of web pages.
- Size of web pages.
- Network **transfer times**.

...and on and on.

Long-tailed distributions



- Pareto distribution \Rightarrow straight line.
- Lognormal distribution \Rightarrow curved line.

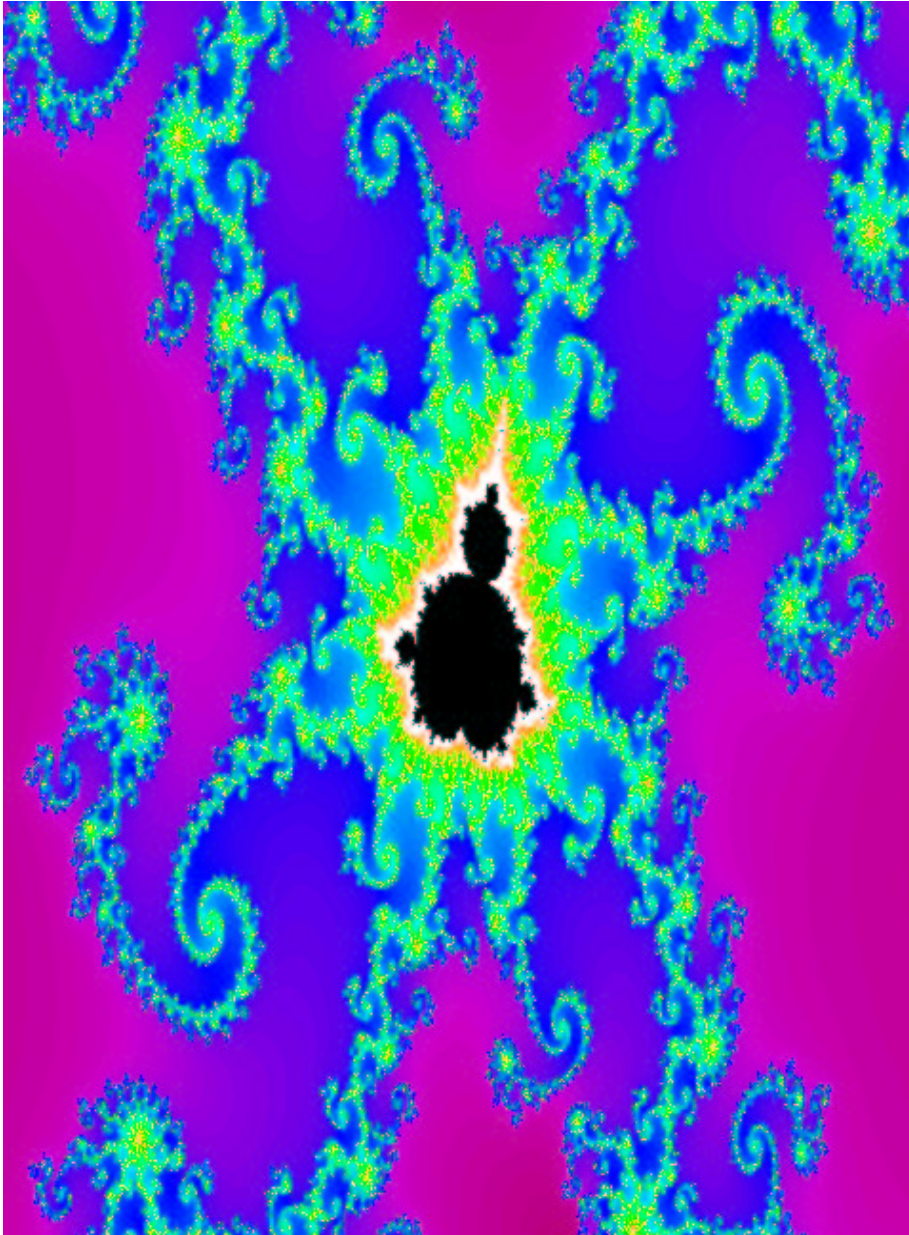
“Evidence for long-tailed distributions in the Internet”, Downey, 2003.

So what?

Long-tailed distributions have several implications:

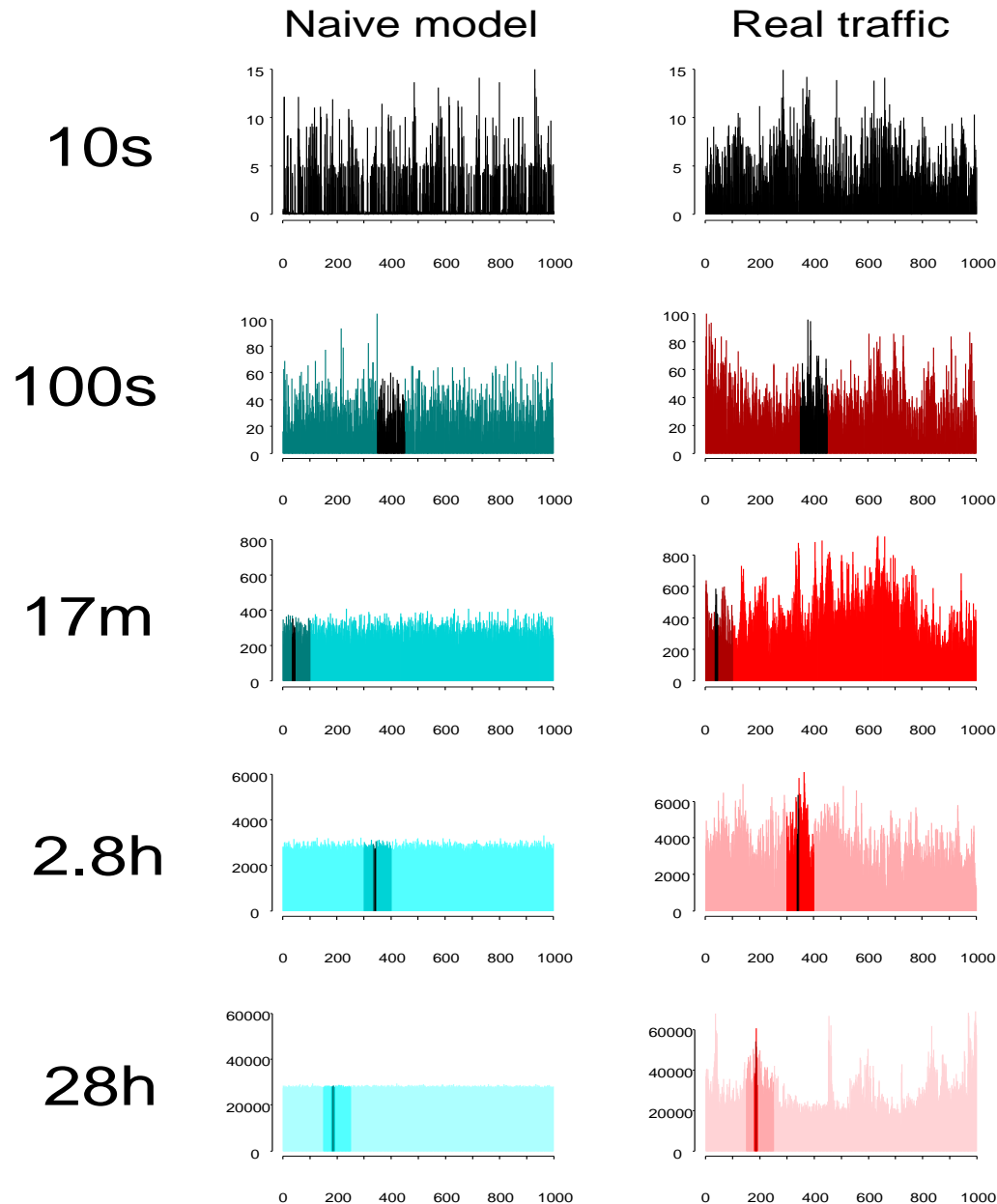
- Remaining lifetime increases with age.
- Network traffic is bursty.
- Network traffic is **self-similar**.

Self-similarity



- Real world: visually similar over a range of spatial scales.
- Fractals: geometrically similar over all spatial scales.
- Time-series: **statistically** similar over a range of **time scales**.

Self-similarity



- Normally, large scales are more **predictable** than small.
- Networks are **bursty** even at large time scales.

“On the Self-Similar Nature of Ethernet Traffic,” Leland, Taqqu, Willinger and Wilson, 1993.

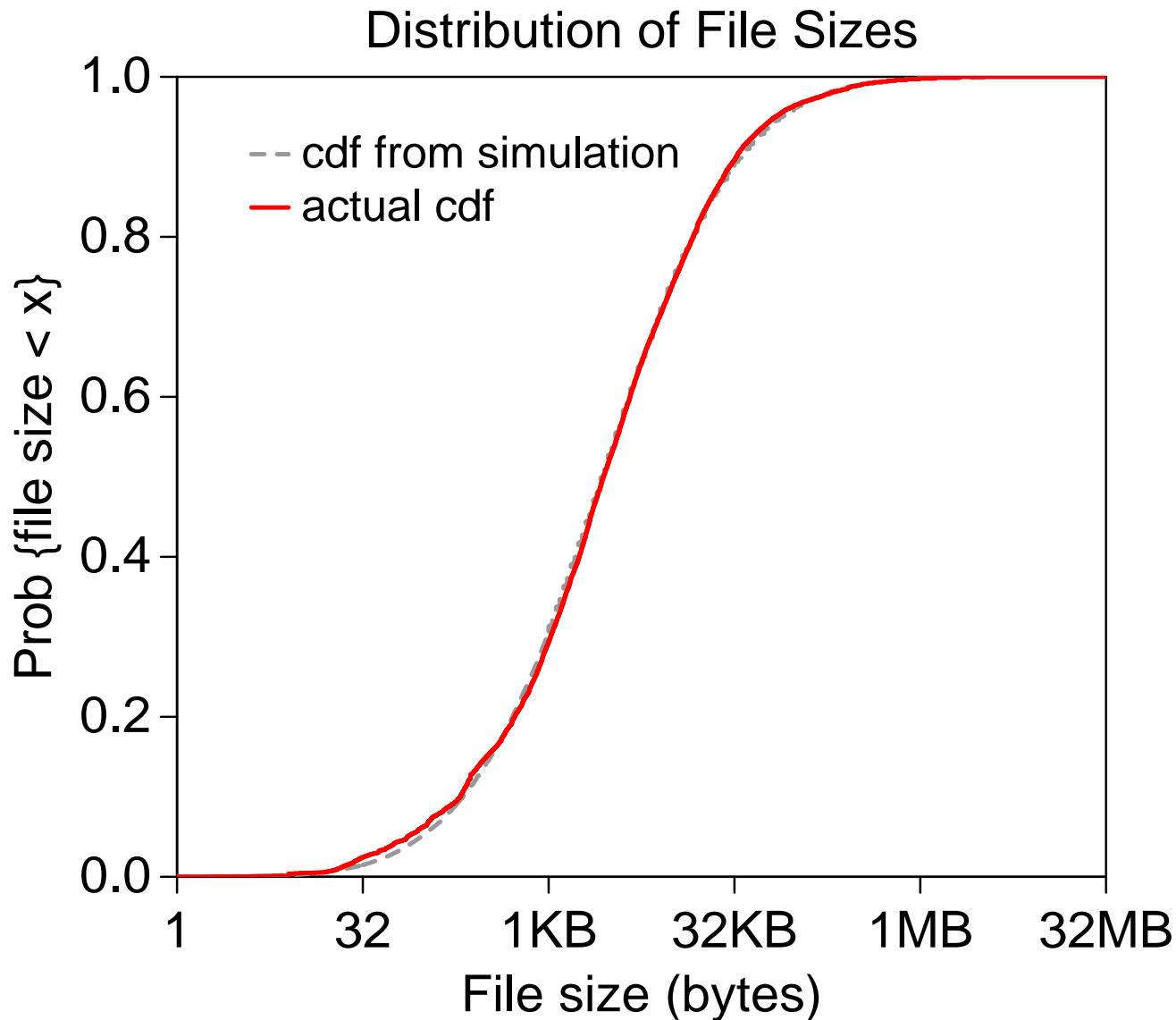
Self-similarity

Best current theory:

- Long-tailed transfer sizes cause self-similarity.

So what causes long-tailed transfer sizes?

File sizes



- File sizes are lognormal.
- Explained by simple model of **user behavior**.

“The structural cause of file size distributions,” Downey, 2001.

User model

Most new files are a

1. copy,
2. modification, or
3. translation

of an existing file.

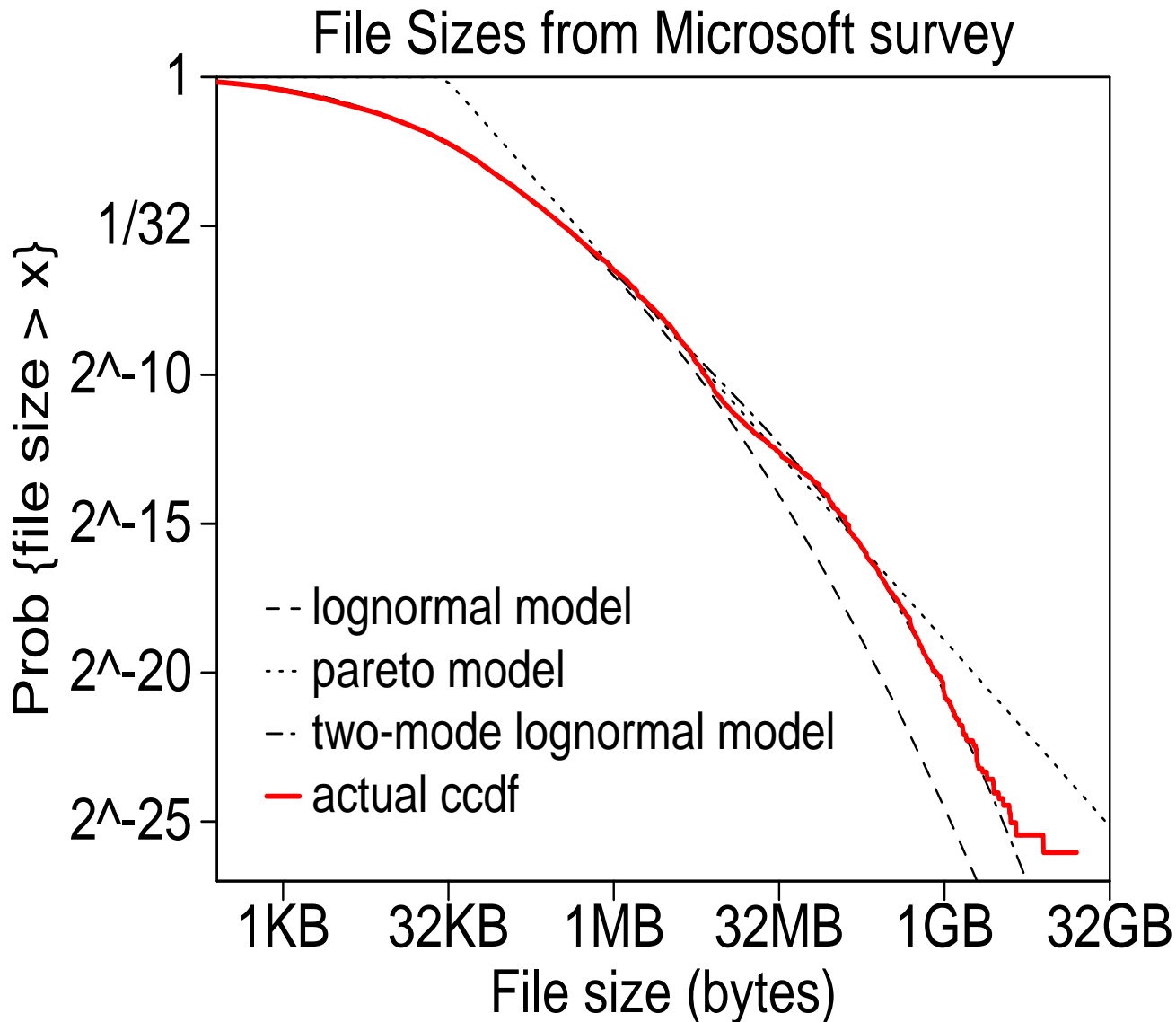
User model

Every day the user:

1. Chooses an **existing file** at random.
2. Chooses a small **multiplier** at random.
3. Creates a **new file** with size = old file size * multiplier.
4. Repeat.

Distribution of file sizes tends toward **lognormal**.

File sizes



- Mixture of lognormal distributions approaches a **Pareto distribution**.
- 40 million files from 10,000 machines.

Self-similarity in the Internet

As we understand it now...

- File systems evolve toward long-tailed size distributions.
- Long tailed-distributions induce **self-similar traffic** patterns.

Self-similarity in the Internet

These patterns are the result of interactions among:

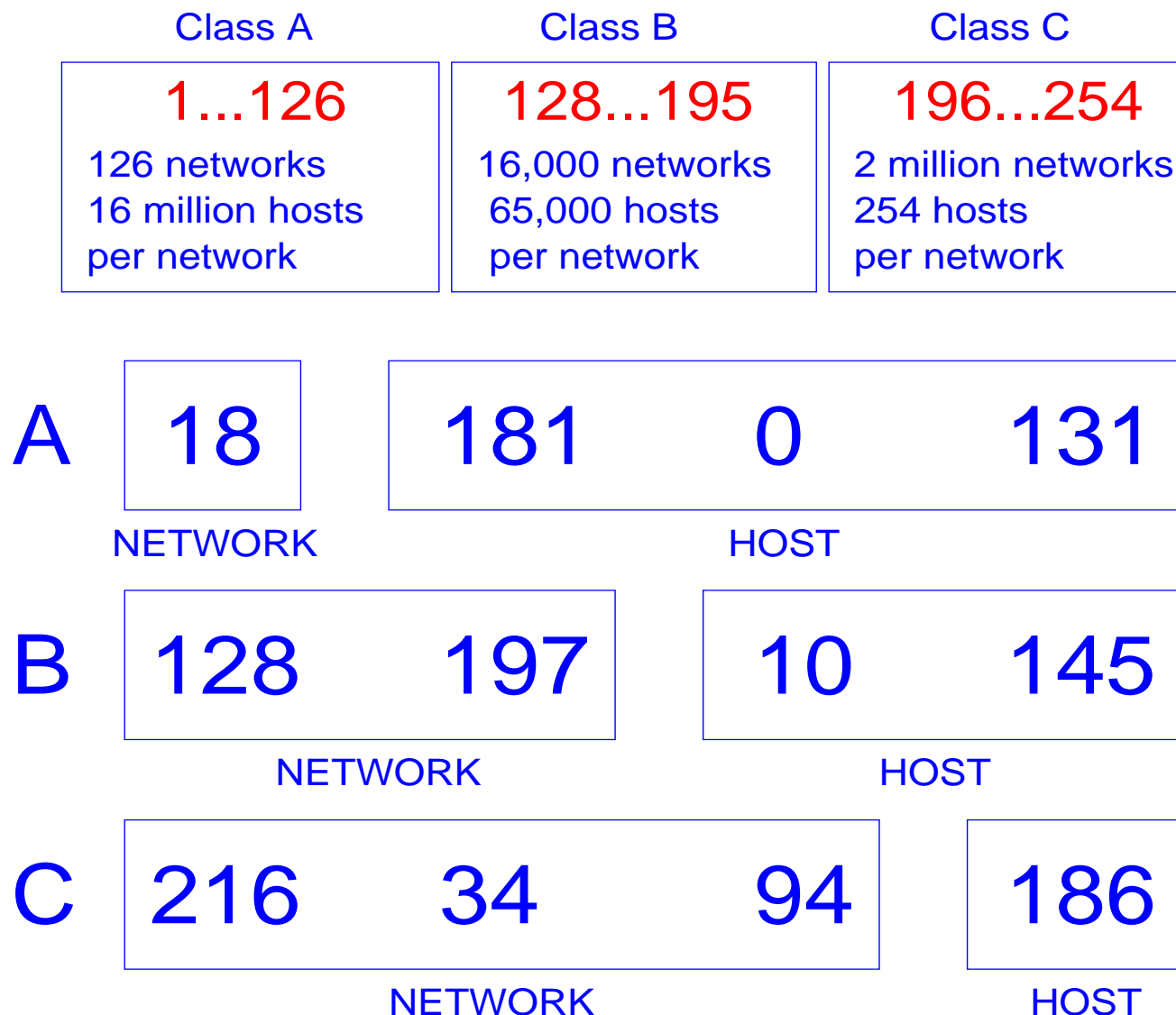
- Technology.
- People.
- Information.

Summary

- Understand networking technology in terms of **layers of abstraction**.
- Each layer solves a set of **problems**, provides a set of **capabilities**.
- Understand the Internet in terms of **underlying technology** and **emergent properties**.

Hierarchical addresses

ICANN hands out blocks of addresses.



Hierarchical addresses

Addresses encode history.

- MIT has a Class A network (18.x.x.x)
- BU has a Class B network (128.197.x.x)
- Jeff Elkner has a Class C network (216.34.94.x)
- How did **Olin** get a Class A address??? (4.x.x.x)

Hierarchical addresses

The pros and cons.

- Inefficient use of addresses (IPv6).
- Common prefix \Rightarrow colocation.

